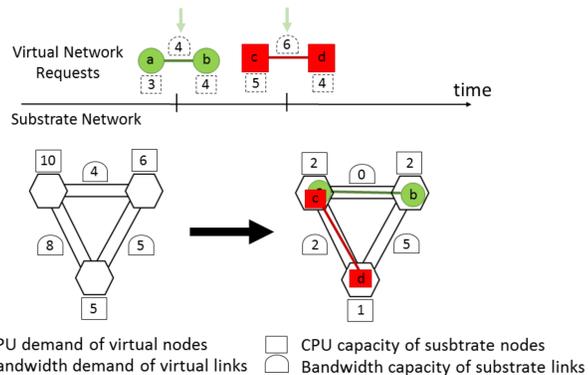


Introduction & Motivation



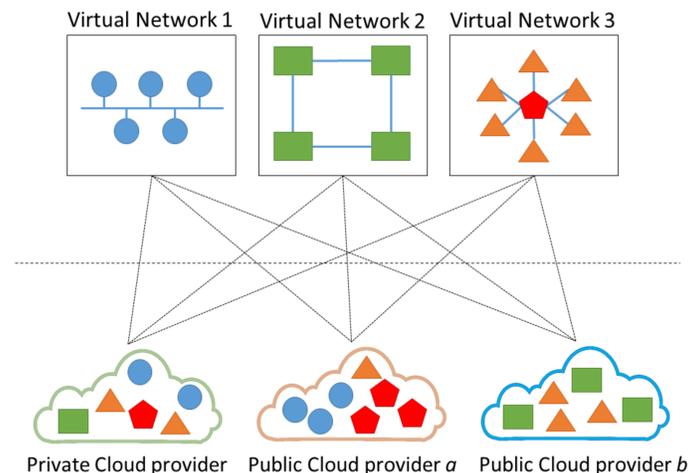
Embedding of virtual networks

In network virtualization, the Virtual Network Embedding (VNE) problem deals with finding an effective mapping of virtual nodes and links onto a sub-

strate network. This problem is traditionally formulated with the objective of maximizing network provider revenue by efficiently embedding incoming virtual network (VN) requests. A mostly unexplored subject on this problem is security [1]. With the advent of **network virtualization platforms** [2], cloud operators now have the ability to extend their cloud computing offerings with VNs. To shift their workloads to the **cloud**, tenants trust their cloud providers to guarantee that their workloads are secure and available. Unfortunately, there is an **increasing number of evidence that problems occur** [3]. Thus, **security and dependability** are becoming critical factors that should be considered by VNE algorithms.

Contribution

We propose a VN embedding solution that considers **security and dependability** as first class citizens. For this purpose, we introduce specific security and dependability constraints. To further extend the resiliency properties of our solution, we **assume a multiple cloud provider model** (e.g., one based on nested virtualization). By not relying on a single cloud provider we avoid internet-scale single points of failures, avoiding cloud outages by replicating workloads across clouds. In addition, we can enhance security by leaving sensitive workloads in the tenant's private clouds.



In a multi-cloud environment, different resources from different clouds can be used by the VNs

MIP Formulation

We formulate the problem as a Mixed Integer Program (MIP). The objective is to minimize the cost of embedding VN requests, taking into account the VN lifetime and the resource's security levels:

$$\min Dur_k^V \left[\sum_{i,j \in N_k^V} \sum_{u,v \in N^S} \lambda f_{u,v}^{i,j} lev^S(u,v) \right]$$

(Reduce the costs of allocated bandwidth)

$$+ \sum_{u,v \in N^S} r_{u,v} lev^S(u,v)$$

(Reduce the costs of allocated bandwidth for backup)

$$+ \sum_{p \in N_k^V} \sum_{w \in N^S} R_N(w) \Theta_{w,p} lev^S(w)$$

(Reduce the costs of allocated CPU)

$$+ \sum_{w \in N^S} \gamma_w lev^S(w)$$

(Reduce the costs of allocated CPU for backup)

We consider the typical flow conservation and resource capacity constraints. We further define security and dependability constraints.

Security constraints:

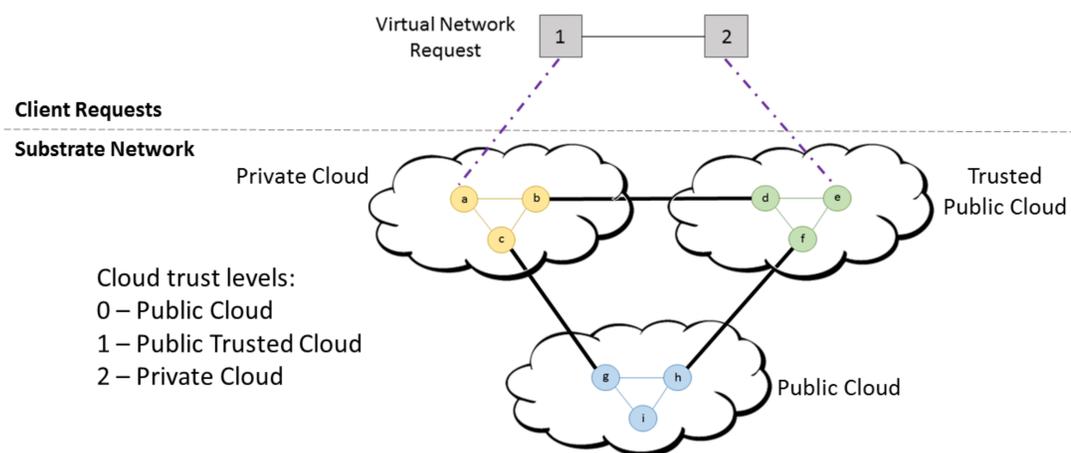
- Security level of a physical node \geq security level required by the virtual node;
- Security level of a physical path \geq security level required by the virtual link;
- Sensitive virtual resources should not be hosted in public clouds.

Dependability Constraints:

- A physical resource should guarantee at least the replication level required by the virtual resource;
- The physical path should guarantee at least the replication level required by the virtual link.

Solution

Security Levels for Virtual Nodes: 0 – Normal container 1 – Normal VM 2 – Secure VM	Security Levels for Virtual Links: 0 – No security 1 – Authenticity & Integrity 2 – Authenticity, Integrity & Confidentiality	Dependability Levels for Virtual Nodes: 0 – No replication 1 – Replication on the same cloud 2 – Replication in other cloud	Dependability Levels for Virtual Links: 0 – No redundancy 1 – Path redundancy
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Security and dependability options for the virtual networks

Different clients/tenants have different security and dependability requirements. As such, in our solution clients specify their requirements as security and dependability levels for each virtual resource (nodes and links). In addition, the client can choose to which specific cloud its virtual resources may be mapped.

References

- [1] A. Fischer et al. Virtual network embedding: A survey. 2013.
- [2] T. Koponen et al. Network virtualization in multitenant datacenters. 2014.
- [3] Cloud Security Alliance. The notorious nine cloud computing top threats in 2013.

Acknowledgments

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Work in progress: heuristic

Due to the complexity inherent to the embedding problem and the considerable size of the problem space in our MIP formulation, this solution is not efficient. Thus, we are currently investigating possible heuristics to this problem.

Our initial idea is to design a two-phase algorithm: the first phase for node mapping, and the second for link mapping. We also consider techniques to correlate the two phases, with node mapping taking into account the subsequent link mapping.

For node mapping we plan to explore a greedy approach that chooses a substrate node according to an utility function, while fulfilling the security and dependability constraints. The utility function takes into account the capacity of each node, namely in terms of CPU resources and availability. The proposed function includes a multiplier factor that gives more value to nodes with more links, to increase the correlation between node mapping and link mapping. For the second phase we will consider both shortest path and multi-commodity flow algorithms (the latter for when the substrate network supports path splitting/multi-path).

As security constraints we include the type of clouds where the resource can be mapped and the security level provided by the substrate nodes and links. For dependability, we consider the use of backup nodes (including replication in different clouds), path redundancy, and an availability factor as part of the utility function.